Asian Journal of Research in Computer Science

3(2): 1-10, 2019; Article no.AJRCOS.47900
ISSN: 2581-8260

Meta-Heuristics Approach to Knapsack Problem in Memory Management

Emmanuel Ofori Oppong[1], Stephen Opoku Oppong[2*], Dominic Asamoah[1] and Nuku Atta Kordzo Abiew [2]

[1]Department of Computer Science, KNUST, Ghana.
[2]Faculty of Computing and Information Systems, GTUC, Ghana.

Authors' contributions

This work was carried out in collaboration among all authors. All authors read and approved the final manuscript.

Article Information

DOI: 10.9734/AJRCOS/2019/v3i230087
Editor(s):
(1) Dr. Jong-Wuu Wu, Professor, Department of Applied Mathematics, National Chiayi University, Taiwan.
Reviewers:
(1) R. Nedunchelian, Anna University, India.
(2) Anthony Spiteri Staines, University of Malta, Malta.
Complete Peer review History: http://www.sdiarticle3.com/review-history/47900

Original Research Article

Received 12 January 2019
Accepted 27 March 2019
Published 06 April 2019

ABSTRACT

The Knapsack Problems are among the simplest integer programs which are NP-hard. Problems in this class are typically concerned with selecting from a set of given items, each with a specified weight and value, a subset of items whose weight sum does not exceed a prescribed capacity and whose value is maximum. The classical 0-1 Knapsack Problem arises when there is one knapsack and one item of each type. This paper considers the application of classical 0-1 knapsack problem with a single constraint to computer memory management. The goal is to achieve higher efficiency with memory management in computer systems.

This study focuses on using simulated annealing and genetic algorithm for the solution of knapsack problems in optimizing computer memory. It is shown that Simulated Annealing performs better than the Genetic Algorithm for large number of processes.

Keywords: Knapsack; memory management; genetic algorithm; simulated annealing.

**Corresponding author: E-mail: sopokuoppong@yahoo.com;*

1. INTRODUCTION

A great variety of practical problems can be represented by a set of entities, each having an associated value, from which one or more subsets has to be selected in such a way that the sum of the values of the selected entities is maximized, and some predefined conditions are respected. The most common condition is obtained by also associating a weight to each entity and establishing that the sum of the entity sizes in each subset does not exceed some prefixed bound. These problems are generally called knapsack problems, since they recall the situation of a traveler having to fill up his knapsack by selecting from among various possible objects those which will give him the maximum comfort. One such problem is in computer memory management.

Modern computer memory management is for some causes a crucial element of assembling current large applications. First, in large applications, space can be a problem and some technology are efficiently needed to return unused space to the program. Secondly, inexpert implementations can result in extremely unproductive programs since memory management takes a momentous portion of total program execution time and finally, memory errors become rampant, such that it is extremely difficult to find programs when accessing freed memory cells. It is much secured to build more unfailing memory management into design even though complicated tools exist for revealing a variety of memory faults. It is for this basis that efficient schemes are needed to manage allocating and freeing of memory by programs.

Optimizing current memory management strategies strength is performed by altering the space allocated to each task. To achieve high levels of multiprogramming while avoiding thrashing such policies vary the load (i.e., the number of active tasks). Additionally, in a system that runs out of capacity probably because the system is undersized, several options are available. This option includes either upgrading the processor (if possible), reduce available functionality, or optimize.

A great deal of realistic problems where some predefined conditions are respected such that the sum of the values of the selected entities is maximized can be represented by a set of entities, each having an associated value, from which one or more subsets has to be selected. The most ordinary situation is obtained by establishing that the sum of the entity sizes in each subset does not exceed some prefixed bound by associating a weight/size to each entity.

The goal of this paper is to maximize the number of processes in a limited memory space.

2. LITERATURE REVIEW

Knapsack problems have been studied intensively in the past decade attracting both theorist and practitioners. The theoretical interest arises mainly from their simple structure which both allows exploitation of a number of combinational properties and permits more complex optimization problems to be solved through a series of knapsack type. From a practical point of view, these problems can model many industrial applications, the most classical applications being capital budgets, cargo loading and cutting stock. In this section a review of literature on knapsack problems and applications is presented.

The knapsack problem (KP) is a traditional combinatorial issue used to show numerous modern circumstances. ―Since Balas and Zemel a dozen years ago introduced the so-called core problem as an efficient way of solving the Knapsack Problem, all the most successful algorithms have been based on this idea. All knapsack Problems belong to the family of NP-hard problems, meaning that it is very unlikely that polynomial algorithms for these problems can be devised [1].

The Knapsack problem has been concentrated on for over a century with prior work dating as far back as 1897. ―It is not known how the name Knapsack originated though the problem was referred to as such in early work of mathematician Tobias Dantzig suggesting that the name could have existed in folklore before mathematical problem has been fully defined [2].

Heuristic algorithms experienced in literature that can generally be named as population heuristics include; ―genetic algorithms, hybrid genetic algorithms, mimetic algorithms, scatter-search algorithms and bionomic algorithms. Among these, Genetic Algorithms have risen as a dominant latest search paradigm [3].

Genetic Algorithms (GA) are PC algorithms that hunt down fine solutions to a problem from among countless solutions. They are versatile heuristic search algorithm in view of the evolutionary thoughts of natural selection and hereditary qualities. “These computational paradigms were inspired by the mechanics of natural evolution, including survival of the fittest, reproduction, and mutation. This algorithm is an intelligent exploitation of random search used in optimisation problems” [4].

Bortfeldt and Gehring presented a hybrid genetic algorithm (GA) for the container packing problem with boxes of unlike sizes and one container for stacking. Generated stowage plans include several vertical layers each containing several boxes. Within the procedure, stowage plans were represented by complex data structures closely related to the problem. To generate offspring, specific genetic operators were used that are based on an integrated greedy heuristic [5].

GAs often calls for the creation and assessment of lots of dissimilar children. However, GAs are capable of generating high-quality solutions to many problems within reasonable computation times [6,7,8,9]. Additionally, while performing search in large state-space or multi-modal state-space, or n-dimensional surface, a genetic algorithm offers significant benefits over many other typical search optimisation techniques like linear programming, heuristic, depth-first, breath-first.

Proposed in [10], simulated annealing maintain a temperature variable to create heating process. The temperature is earlier set high and after that allows to gradually "cool" as the algorithm runs. While this temperature variable is high the algorithm will be permitted, with more recurrence, to accept solutions that are more awful than the present solution. This gives the algorithm the capacity to hop out of any local optimums it discovers itself on early on in execution. As the temperature is decreased so is the possibility of tolerating more awful solution, thus permitting the algorithm gradually focusing on a zone of the search space in which ideally, a near ideal solution can be found.

Simoes and Costa [11] performed an empirical study and evaluated the exploits of the transposition A-based Genetic Algorithm (GA) and the classical GA for solving the 0/1 knapsack problem. Obtained results showed that, just like in the domain of the function optimization, transposition is always superior to crossover.

Eager about making use of a easy heuristic scheme (simple flip) for answering the knapsack problems, [12] offered a study work on the application of usual zero-1 knapsack trouble with a single limitation to determination of television ads at significant time such as prime time news, news adjacencies, breaking news and peak times.

Martello et al. [13] presented a new algorithm for the optimal solution of the 0-1 Knapsack problem, which is particularly effective for large-size problems. The algorithm is based on determination of an appropriate small subset of items and the solution of the corresponding "core problem": from this they derived a heuristic solution for the original problem which, with high probability, can be proved to be optimal. The algorithm incorporated a new method of computation of upper bounds and efficient implementations of reduction procedures.

Huttler and Mastrolilli [14] addressed the classical knapsack problem and a variant in which an upper bound is imposed on the number of items that can be selected. It was shown that appropriate combinations of rounding techniques yield novel and more powerful ways of rounding. Moreover, they presented a linear-storage polynomial time approximation scheme (PTAS) and a fully polynomial time approximation scheme (FPTAS) that compute an approximate solution, of any fixed accuracy, in linear time. These linear complexity bounds give a substantial improvement of the best previously known polynomial bounds.

Hanafi and Freville [15] described a new approach to tabu search (TS) based on strategic oscillation and surrogate constraint information that provides a balance between intensification and diversification strategies. New rules needed to control the oscillation process are given for the 0 /1 multidimensional knapsack (0/1 MKP). Based on a portfolio of test problems from the literature, our method obtains solutions whose quality is at least as good as the best solutions obtained by previous methods, especially with large scale instances. These encouraging results confirm the efficiency of the tunneling concept coupled with surrogate information when resource constraints are present.

Rinnooy et al. [16] proposed a class of generalized greedy algorithms is for the solution of the multi-knapsack problem. Items are selected according to decreasing ratios of their profit and a weighted sum of their requirement coefficients. The solution obtained depended on the choice of the weights. A geometrical representation of the method was given and the relation to the dual of the linear programming relaxation of multi-knapsack is exploited. They investigated the complexity of computing a set of weights that gives the maximum greedy solution value. Finally, the heuristics were subjected to both a worst-case and a probabilistic performance analysis.

Balachandar and Kannan [17] presented a heuristic to solve the 0/1 multi-constrained knapsack problem (0/1 MKP) which is NP-hard. In this heuristic the dominance property of the constraints is exploited to reduce the search space to find near optimal solutions of 0/1 MKP. This heuristic was tested for 10 benchmark problems of sizes up to 105 and for seven classical problems of sizes up to 500, taken from the literature and the results were compared with optimum solutions. Space and computational complexity of solving 0/1 MKP using this approach were also presented. The encouraging results especially for relatively large size test problems indicate that this heuristic can successfully be used for finding good solutions for highly constrained NP-hard problems.

Elhedhli [18] considered a class of nonlinear knapsack problems with applications in service systems design and facility location problems with congestion. They provided two linearizations and their respective solution approaches. The first is solved directly using a commercial solver. The second is a piecewise linearization that is solved by a cutting plane method.

Devyaterikova et al. [19] presented discrete production planning problem which may be formulated as the multidimensional knapsack problem is considered, while resource quantities of the problem are supposed to be given as intervals. An approach for solving this problem based on using its relaxation set is suggested. Some *L*-class enumeration algorithms for the problem are described. Results of computational experiments were presented.

Chen et al. [20] presented pipeline architectures for the dynamic programming algorithms for the knapsack problems. They enabled them to achieve an optimal speedup using processor arrays, queues, and memory modules. The processor arrays can be regarded as pipelines where the dynamic programming algorithms are implemented through pipelining.

3. METHODOLOGY

Because of their wide range of applicability, knapsack problems have known a large number of variations such as: single and multiple-constrained knapsacks, knapsacks with disjunctive constraints, multidimensional knapsacks, multiple choice knapsacks, single and multiple objective knapsacks, integer, linear, non-linear knapsacks, deterministic and stochastic knapsacks, knapsacks with convex / concave objective functions, etc.

This is a 0-1 knapsack problem, pure integer programming with single constraint which forms a very important class of integer programming.

The 0-1 Knapsack Problem (KP) can be mathematically formulated through the following integer linear programming [21].

$$\text{Maximize} \sum_{j=1}^{n} P_j x_j \tag{1}$$

$$\text{Subject to} = \sum_{j=1}^{n} (w_j x_j) \le c \tag{2}$$

$$x_j = 0 \text{ or } 1,\ j = 1, \ldots, n$$

Where, P_j refers to the value, or worth of item j, x_j refers to the item j, w_j refers to the relative-weight of item *j*, with respect to the knapsack and C refers to the capacity, or weight-constraint of the knapsack. There exist *j* = 1,…,*n* items, and there is only one knapsack.

The use of two major meta-heuristics approaches, Genetic algorithm and Simulated annealing which have been used to solve large scale problems [22] will be considered in this paper.

3.1 Simulated Annealing

Simulated annealing (SA) is a local search algorithm capable of escaping from local optima. Its case of implementation, convergence properties and its capability of escaping from local optima has made it a popular algorithm over the past decades. Simulated annealing is so named because of its analogy to the process of physical annealing with solids in which a crystalline solid is heated and then allowed to cool very slowly until it achieves stable state. i.e. its minimum lattice energy state and thus is free of crystal effects. Simulated annealing mimics this type of thermodynamic behavior in searching for global optima for discrete optimization problems (DOP) [23].

To formally describe simulated annealing algorithm for KP, some definitions are needed. Let Ω be the solution space: Define η(ω) to be the neighborhood function for w ∈ Ω. Simulated annealing starts with an initial solution ω ∈ Ω. A neighborhood solution $\omega^1 \in \eta(\omega)$ is then generated randomly in most cases. Simulated annealing is based on the Metropolis acceptance criterion, which models how a thermodynamic system moves from its current solution ω ∈ Ω to a candidate solution $\omega i \in \eta(\omega)$ in which the energy content is being minimized. The candidate solution ω^1 is accepted as the current solution based on the acceptance probability.
In this survey, finite-time implementations of simulated annealing algorithm are considered, which can no longer guarantee to find an optimal solution, but may result in faster executions without losing too much on the solution quality. Simulated annealing algorithm with static cooling schedule [24] for KP is outlined in pseudo-code.

```
1  Select an initial solution $\omega = (x_1,\ldots, x_n) \in \Omega$; an initial temperature $t = t_0$;
2  control parameter value α ; final temperature e; a repetition schedule, M that defines the number of iterations executed at each temperature;
3  $\text{Incumbent solution} \leftarrow f(\omega)$;
4  Repeat;
5  Set repetition counter $m = 0$;
6  Repeat;
7  Select an integer i from the set $\{1,2,\ldots,n\}$ randomly.
8  If $x_i = 0$, pick up item i, i.e. set $x_i = 1$, obtain new solution ω1 then
9  while solution ω1 is infeasible, do
10 drop another item from ω randomly; denote the new solution as ω1
11 let $\Delta = f(\omega 1) - f(\omega)$
12 while $\Delta \geq 0$ or Random (0,1) < $e^{\Delta/t}$ do $\omega \leftarrow \omega 1$
13 Else
14 drop item i and pick another item randomly, get new solution ω1
15 let $\Delta = f(\omega 1) - f(\omega)$
16 while $\Delta \geq 0$ or Random (0,1) < $e^{\Delta/t}$ do $\omega \leftarrow \omega 1$
17 End If
18 If incumbent solution < $f(\omega)$, Incumbent solution $\leftarrow f(\omega)$
19 $m = m + 1$;
20 Until $m = M$
21 set $t = a * t$;
22 Until $t < e$
```

A set of parameters needs to be specified that govern the convergence of the algorithm, i.e. initial temperature to , temperature control parameter α, final temperature e, and Markov chain length M, in order to study the finite-time performance of simulated annealing algorithm. Here t_o should be the maximal difference in cost between any two neighboring solutions [24].

The parameters used for the Simulated Annealing are:

Cooling factor:	0.98
Termination Temperature:	0.2
Initial Temperature:	100
Neighbor Sampling Size:	350

3.2 Genetic Algorithm

A genetic algorithm (GA) can be described as an "intelligent" probabilistic search algorithm and is based on the evolutionary process of biological organisms in nature. During the course of evolution, natural populations evolve according to the principles of nature selection and "survival of the fittest." Individuals who are most successful in adapting to their environment will have a better chance of surviving and reproducing, while individuals who are less fit will be eliminated. This means that the genes from highly fit individuals will spread to an increasing number of individuals in each successive generation. The combination of good characteristics from highly adapted parents may produce even more fit offspring. In this way, species evolve to become increasingly better adapted to the environment [25].

A GA simulates these processes by taking an initial population of individuals and applying genetic operators in each reproduction. In optimization terms, each individual in the population is encoded into a string or chromosome that represents a possible solution to a given problem. The fitness of an individual is evaluated with respect to a given objective function. Highly fit individuals or solutions are given opportunities to reproduce by exchanging pieces of their genetic information in a crossover procedure with other highly fit individuals. This produces new "offspring" solutions (i.e. children) who share some characteristics taken from both parents. Mutation is often applied after crossover by altering some genes in the strings. The offspring can either replace the whole population (generational approach) or replace fewer fit individuals (steady-state approach). This evaluation-selection-reproduction cycle is repeated until a satisfactory solution is found.

The basic steps of a simple GA are shown below [26]

Step 1: Generate an initial population
Step 2: Evaluate fitness of individuals in the population

The objective function value ($\sum_{j=1}^{n} pjXj$) equates to how good a solution is, that is, its fitness.

In general, an initial population is randomly generated in some way.

Step 3: Repeat

a. Select individuals from the population to be parents
   In the GA world for the KP, parents will be chosen by binary tournament selection. In binary tournament selection, two individuals are randomly selected from the population. From these two, the individual with the best fitness is selected to be the first parent
b. Recombine (mate) parents to produce children
   In the GA world for the KP, a single child will be obtained from two parents by uniform crossover. In uniform crossover each bit in the child solution is created by: Repeat for each bit in turn choose one of the two parents at random set the child bit equal to the bit in the chosen parent in one-point crossover, a pint between two adjacent bits is randomly selected, “cut” the parents into two segments and create two children by rejoining the segments.
c. Mutate the children Evaluate fitness of the children
   Mutation corresponds to small changes that are stochastically applied to the children mutation can be applied with a constant probability or with an adaptive probability that changes over the course of the algorithm (perhaps in response to the number of iterations that have passed or in response to population characteristics).
d. Replace some or all of the population by the children until

Step 4: You decide to stop whereupon report the best solution encountered

The parameters used for the Genetic Algorithm are:

Population Size: 500
Recombination Rate:0.7
Mutation Rate: 0.005
Number of Crossover Points: 3

3.3 Chi-square

To ascertain whether the time taken and memory sued to obtain a solution is dependent or not on the number of processes, the chi-square test is used. The chi-square test of independence is a statistical test to determine if two or more classifications of the samples are independent or not.

The chi-square test is computed with the following equation [27]

$$\chi^2 = \sum_{i}^{k} \frac{(O_i - E_i)^2}{E_i} \tag{3}$$

Where:

O_i is the cell frequencies actually observed in category i
E_i is the cell frequencies that would be expected in category i, if the two tables were statistically independent
k is the total number of cells or categories

Hypothesis testing is a statistical method that is used in making statistical decisions using experimental data. Hypothesis tests are used to test the validity of a claim that is made about a population (in this context all data under consideration). This claim made, in essence, is called the null hypothesis (H_0) and the alternative hypothesis (H_1) is the one considered if the null hypothesis is concluded to be untrue. For this paper, the hypothesis is to ascertain whether the time taken or the memory used to obtain a solution is dependent on the number of processes. Therefore, the null hypothesis (H_0) states that no association exists between the two variables i.e. the variables are statistically independent and the alternate hypothesis (H_1) states that the two variables are related. In performing a hypothesis test in statistics, a p-value helps determine the significance of the results. The p-value can be estimated using the chi-square distribution table or using a statistical package. Before the hypothesis test is performed, a threshold value is chosen, called the significance level of the test and is denoted by α.

The p-value is a number always between 0 and 1 and interpreted in the following way:

- A small p-value (≤ 0.05) indicates strong evidence against the null hypothesis, so you reject the null hypothesis if the significance level $\alpha(=0.05)$.

- A large p-value (> 0.05) indicates weak evidence against the null hypothesis, so you fail to reject the null hypothesis if the significance level $\alpha(=0.05)$.

R (a statistical package) is used to calculate the chi-square value and p-value using this pseudocode.

```
x<=matrix(data)
View(x)
Chisq.test(x)
```

4. ANALYSIS AND RESULTS

Category A: The computer system with a total of 10 created processes, all with their system information in figures. The computer memory can accommodate capacity of 50mb but the total memory of the process is 56 with a combined process activity (number of times process is accessed of 123.

Table 1. Results for category A

	GA	SA
No. of Processes Used	9	9
Memory Used	46	46
Number of Times Process Is Accessed	119	119

From Table 1, it could be seen that all three algorithms provide the same output in terms of all the parameters under consideration. This means that both DP, GA and SA.

Category B: The table below shows a computer system with a total of 50 created processes, all with their system information in figures. The computer memory can accommodate capacity of 100mb. but the total memory of the process is 281 with a combined process activity (number of times process is accessed of 483.

Table 2. Results for category B

	GA	SA
No. of Processes Used	25	23
Memory Used	100	100
Number of Times Process Is Accessed	327	328

From Table 2, GA provided a slight advantage of in terms of the number of process used. Apart from that all three algorithms provided fairly the same result.

Category C: The table below shows a computer system with a total of 100 created processes, all with their system information in figures. The computer memory can accommodate capacity of 300mb. but the total memory of the process is 574 with a combined process activity (number of times process is accessed of 1011.

Table 3. Results for category C

	GA	SA
No. of Processes Used	61	62
Memory Used	300	300
Number of Times Process Is Accessed	815	803

Table 3 shows that DP provides a better result than the rest. All memory needed was utilized showing efficient use of memory available.

Category D: The table below shows a computer system with a total of 500 created processes, all with their system information in figures. The computer memory can accommodate capacity of 1000mb. but the total memory of the process is 2661 with a combined process activity (number of times process is accessed of 5287.

Table 4. Results for category D

	GA	SA
No. of processes used	258	252
Memory used	1000	1000
Number of times process is accessed	3551	3431

Category E: The table below shows a computer system with a total of 1000 created processes, all with their system information in figures. The computer memory can accommodate capacity of 5000mb. but the total memory of the process is 5626 with a combined process activity (number of times process is accessed of 10480).

Table 5. Results for category E

	GA	SA
No. of Processes Used	915	916
Memory Used	5000	5000
Number of Times Process Is Accessed	10299	10307

GA and SA provide fairly the same results in Tables 4 and 5.

The main criteria in evaluating the efficiency of an algorithm is time and space. Even though in terms of results the three algorithms provided similar results, their efficiency will be determined based on the time it took to produce the results and the amount of memory resource it took on the computer.

Table 6. Results for based on time taken

	TIME (ms)	
No. of process	**GA**	**SA**
10	436	60
50	323	52
100	385	87
500	1374	300
1000	2338	554

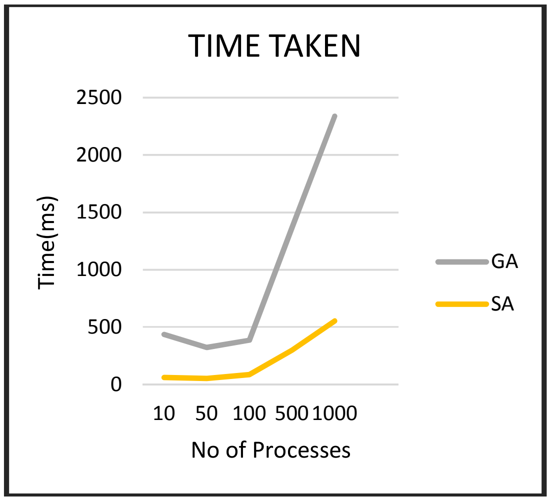

Fig. 1. Results for based on time taken

From Table 6 and Fig. 1, It is seen that GA took more time in giving an optimum out than SA for larger number of processes. As the number of processes increases, time taken increases exponentially for GA as compared to SA.

Also the GA also used more memory utilization for than SA from Table 7 and Fig. 2. The GA outperformed the Sa only when the number of processes.

Using the chi-square test on Table 6, the null and alternate hypothesis are defined as follows:

H_0: Time taken to obtain a solution is independent of number of processes.
H_1: Time taken to obtain a solution is not independent of number of processes.

The chi-square statistic (χ^2)= 18.7547.
The p-value is 0.000878.

Since the p-value of 0.000878 is less than the significance level of 0.05, the null hypothesis(H_0) which stated that the time taken to obtain a solution is independent of number of processes is rejected. This implies that number of processes is dependent on the time taken to obtain a solution

Table 7. Results for based on memory taken

	Memory (byte)	
No. of process	**GA**	**SA**
10	28880312	42511800
50	92815928	45555312
100	100774992	73927720
500	210273904	117057112
1000	233449048	210256440

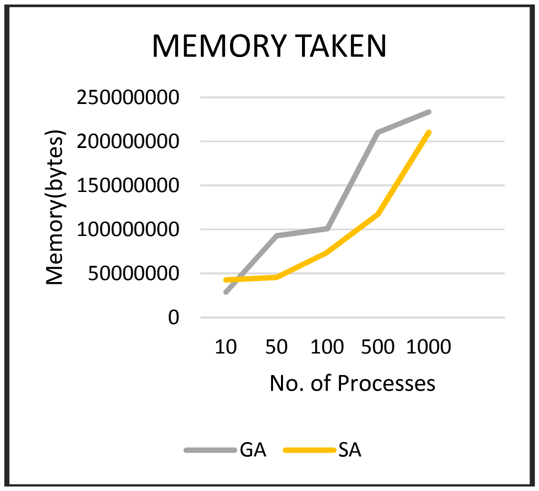

Fig. 2. Results for based on memory taken

Using the chi-square test on Table 7, the null and alternate hypothesis are also defined as follows:

H_0: Memory used in obtaining a solution is independent of number of processes.
H_1: Memory used in obtaining a solution is not independent of number of processes.

The chi-square statistic (χ^2)= 22.8798
The p-value is 0.000134.

Since the p-value of 0.000134 is less than the significance level of 0.05, the null hypothesis is also rejected here. This implies that memory used to obtain a solution is dependent on the number of processes.

5. CONCLUSION AND RECOMMENDATIONS

This paper showed that memory optimization as well as knapsack problem can be successfully solved using heuristic algorithms. In this paper, meta-heuristic algorithms i.e. simulated annealing and genetic algorithm were testes compared for their efficiency in optimizing memory. From Fig. 2, it can be seen that with increase in number of processes, experiments with simulated annealing gives better result than the Genetic Algorithm in terms of both time-taken to obtain a solution and memory taken. From the analysis, it can be seen that for smaller number of processes the GA and SA performance are identical but as the number of processes increases, SA performs better than GA. Therefore, it is concluded that, the most efficient algorithm in knapsack optimizing among the two for large number of processes is Simulated Annealing.

Notwithstanding it extensive use, both SA and GA have their limitations. For SA, If the starting temperature is very high, the search will be a random local search for a period of time i.e. accepting all neighbors during the initial phase of the algorithm. Also, In the SA algorithm, the temperature is decreased gradually. If the temperature is decreased slowly, better solutions are obtained but with a more significant computation time. For GA, if reproduction fails to produce good chromosomes then convergence in the right direction is not possible.

COMPETING INTERESTS

Authors have declared that no competing interests exist.

REFERENCES

1. Pisinger D. Core problems in knapsack algorithms. Operations Research. 1999; 47(4):570-575.
2. Kellerer H, Pferschy U, Pisinger D. Knapsack problems. Springer, Berlin Heidelberg; 2004.
3. Chu PC, Beasley JE. A genetic algorithm for multidimensional knapsack problem. Journal of Heuristics. 1998;4(4):63-68.
4. Sinapova L. An introduction to algorithms. Simpson College, Department of Computer Science 701 North C Street, Indianola IA 50125; 2014.
5. Bortfeldt A, Gehring H. A hybrid genetic algorithm for the container loading problem [J]. European Journal of Operational Research. 2001;131(1):143-161.
6. Beasley JE, Chu PC. A genetic algorithm for the set covering problem. European Journal of Operations Research. 1996; 94(2):392-404.
7. Chu PC, Beasley JE. A genetic algorithm for generalized assignment problem. Computer and Operations Research 1997; 24(1):17-23.
8. Chu PC, Beasley JE. A genetic algorithm for multidimensional knapsack problem. Journal of Heuristics. 1998;4(1):63-86.
9. Chang JT, Meade N, Beasley JE, Sharaiha YM. Heuristics for cardinality constrained portfolio optimization. Computer and Operations Research. 1999;27(2000): 1271-1302.
10. Kirkpatrick S, Gelett CD, Vecchi MP. Optimization by simulated annealing. Science. 1983;621-630.
11. Simoes A, Costa E. Using genetic algorithm with asexual transposition. Proceedings of the genetic and evolutional computation conference. 2001;323-330.
12. Amponsah SK, Oppong EO, Agyeman E. Optimal television adverts selection, case study: Ghana Television (GTV). Research Journal of Information Technology. 2011; 3(1):49-54.
13. Martello S, Pisinger D, Toth P. New trends in exact algorithms for the 0–1 knapsack problem. European Journal of Operations Research. 2000;123:325-332.
14. Mastrolilli M, Huttler M. Hybrid rounding techniques for knapsack problems. Journal of discrete applied mathematics. 2006; 154(4):640-649.
15. Hanafi S, Freville A. An efficient tabu search approach for the 0–1 multi-dimensional knapsack problem. European Journal of Operations Research. 1998; 106(2-3):659-675.
16. Rinnooy K, Stougie AHGL, Vercellis C. A class of generalized greedy algorithms for the multi-knapsack problem. European Journal of Operation Research; 1993.
17. Balachandar R, Kannan K. A new polynomial time algorithm for 0–1 multiple knapsack problem based on dominant principles; 2008. Available:www.sciencedirect.com
18. Elhedhli S. Exact solution of a class of nonlinear knapsack problems. Operations Research Letters. 2005;33(6):615-624.
19. Devyaterikova MV, Kolokolov AA, Kolosov AP. L-class enumeration algorithms for a discrete production planning problem with

interval resource quantities. Computers & Operations Research. 2009;36(2):316-324.

20. Chen G, Maw-Sheng Chern, Jin-Hwang Jang. Pipeline architectures for dynamic programming algorithms; 1990.
Available:www.sciencedirect.com
21. Hisatoshi S. A generalized knapsack problem with variable coefficients, mathematical programming. North-Holland Publishing Company. 1975;162-176.
22. Asamoah D, Baidoo E, Oppong S, Optimizing memory using knapsack algorithm. International Journal of Modern Education and Computer Science (IJMECS). 2017;9(5):34-42.
23. Fubin Q, Rui D. Simulated annealing for the 0/1 multidimensional knapsack problem, numerical mathematics. A Journal of Chinese Universities. 2007; 16(4):320-327.
24. Oppong OE. Optimal resource allocation using knapsack problems: A case study of television advertisements at GTV. Master's Thesis, KNUST; 2009.
25. Djannaty F, Doostdar S. A hybrid genetic algorithm for the multidimensional knapsack problem. International Journal of Contemp. Math. Sciences. 2008;3(9):443–456.
26. Carr J. An introduction to genetic algorithms; 2014.
27. McHugh ML. The chi-square test of independence. Biochemia Medica. 2013; 23(2):143-149.

Peer-review history:
The peer review history for this paper can be accessed here:
http://www.sdiarticle3.com/review-history/47900